

## Unit 2 Boolean Algebra

### 2.1 Introduction

We will use variables like  $x$  or  $y$  to represent inputs and outputs (I/O) of a switching circuit. Since most switching circuits are 2 state devices (having only 2 different possible values) the symbols “0” & “1” are used for the value

Thus for variable  $x$ ,  $x = 0$  or  $x = 1$

Note: Although the symbols 0 & 1 *look like* binary numbers, they are not. They have no numeric value! They merely represent the 2 possible states of a switching circuit.

### 2.2 Basic Operations

Complement: Often called inverse; designated by a '(prime). Complement of 1 is 0 and of 0 is 1

Circuit Representation

Often called NOT operation because  $x = 1$  if  $x$  is not equal to 0

AND: Referred to as a Boolean multiplier, Designated by  $\bullet$

Definition:

$$0 \bullet 0 = 0 \qquad 0 \bullet 1 = 0 \qquad 1 \bullet 0 = 0 \qquad 1 \bullet 1 = 1$$

Note: 0 & 1 are Boolean constants not binary #'s

For  $C = A \bullet B$

A	B	C = A * B
0	0	0
0	1	0
1	0	0
1	1	1

Note:  $C = 1$  if and only if (iff) A and B are both 1

Circuit Representation:

Note: Often the “•” is omitted leaving  $AB$  instead of  $A \bullet B$

OR: Referred to as logic addition, Described by “+”

Definition:

$$0 + 0 = 0 \quad 0 + 1 = 1 \quad 1 + 0 = 1 \quad 1 + 1 = 1$$

For  $C = A + B$

A	B	C = A + B
0	0	0
0	1	1
1	0	1
1	1	1

Note:  $C = 1$  iff A or B (or both) = 1

Circuit Representation

Switches:

$X = 0$  if switch is open

$x = 1$  if switch is closed

Series switches

$T = 0$  open circuit between 1 and 2

$T = 1$  closed circuit between 1 and 2

$T = 1$  iff  $A$  and  $B$  are closed

$$T = A \bullet B$$

### Parallel switches

$T = 1$  iff  $A$  or  $B$  is closed

$$T = A + B$$

Thus series switch perform AND parallel OR operations

## **2.3 Boolean Expressions and Truth Tables**

Example Boolean expressions

$$AB' + C$$

$$[A(C+D)]' + BE$$

Order of operations

*Parentheses*

*Complementation*

*AND*

*OR*

Circuits for above equations

The expression can evaluate by substituting 0 or 1 as values for each variable

Literal: each occurrence of a variable or the complement of an expression

Truth Tables: (table of combinations) Specify values of a Boolean expression for every possible combination of values of the variables in the expression

Ex  $F = A' + B$

A	B	A'	F = A' + B
0	0	1	1
0	1	1	1
1	0	0	0
1	1	0	1

An expression with n variables, each having a value of 0 or 1 has a different number of combinations of values of variable =  $2^n$

See Table 2-1 for  $AB' + C$  Example

## 2.4 Basic Theorems

$$X + 0 = X$$

$$X \bullet 0 = 0$$

$$X + 1 = 1$$

$$X \bullet 1 = X$$

Idempotent laws

$$X + X = X$$

$$X \bullet X = X$$

Involution law

$$(X')' = X$$

Laws of complementary

$$X + X' = 1 \qquad X \bullet X' = 0$$

Any of these theorems can be proven by substituting in a value of 0 or 1 for X

Note: An entire expression can be substituted for x

Switch examples

$$A * A$$

$$A + A$$

$$A + 1$$

$$A + 0$$

## **2.5 Commutative, Associate, Distributive Laws**

### Commutative Laws

$$XY = YX \qquad X + Y = Y + X$$

### Associative Laws

$$(XY)Z = X(YZ) = XYZ$$

$$(X + Y) + Z = X + (Y + Z) = X + Y + Z$$

Proof using truth table:

<u>X</u>	<u>Y</u>	<u>Z</u>	<u>XY</u>	<u>YZ</u>	<u>(XY)Z</u>	<u>X(YZ)</u>
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	1	0	0
1	0	0	0	0	0	0
1	0	1	0	0	0	0
1	1	0	1	0	0	0
1	1	1	1	1	1	1

### Circuit Gates

When ANDing variables the result is a 1 iff all values are 1

When OR the result is a 1 if any value is a 1 and 0 iff all variables are zero

### Distributive Laws

$$X(Y+Z) = XY + XZ$$

$$X + YZ = (X + Y)(X + Z)$$

Proof

This law is a very useful Boolean manipulation

## 2.6 Simplification Theorems

$$XY + XY' = X \quad (X+Y)(X+Y') = X$$

$$X + XY = X \quad X(X + Y) = X$$

$$(X + Y')Y = XY \quad XY' + Y = X + Y$$

Each can be proven with truth tables or algebraically

$$\begin{aligned} \text{Ex. } X + XY &= X \bullet 1 + X \bullet Y \\ &= X(1 + Y) \\ &= X \end{aligned}$$

Ex.  $Y + XY'$  using switches

Ex.  $F = A(A' + B)$  using gates

Example Problems:

$$Z = A'BC + A'$$

$$Z = [A + B'C + D + EF] [A + B'C + (D + EF)']$$

$$Z = (AB + C) (B'D + C'E') + (AB + C)'$$

## 2.7 Multiplying Out, Factoring

Multiplying

Sum of products- when all products are products of only single variables  
End result when an expression is fully multiplied out

$$\begin{array}{l} \text{Ex. } AB + CD + EF \\ \text{or} \\ A + BC + DE \end{array}$$

*Distributive laws are used to obtain the sum of products*

Note: When multiplying out try second distributive law first, if possible

$$\text{Ex. } (A + BC) (A + D + E)$$



The problem could be done by multiplying out fully but that would be time consuming

$$(A + BC)(A + D + E)$$

Factoring

Product of sums- *all sums are sums of single variables*

*An expression is fully factored iff it is in its product of sum form*

Ex.  $A + B'CD$

Ex.  $AB' + C'D$

Ex.  $C'D + C'E' + G'H$

Note: *When factoring apply the ordinary distributive law before the second law*

## Gates

Sum of products - expressed as one or more AND gates feeding a single OR gate

$$\text{Ex. } A + BCD + BCE$$

Product of sums - expressed as one or more OR gates feeding a single AND gate

$$\text{Ex. } (A + C)(A + B)(A + D)$$

## **2.8 DeMorgan's Law**

Used to find inverse or complement of Boolean expressions

$$(X + Y)' = X'Y'$$

$$(XY)' = X' + Y'$$

Proof:

X	Y	X'	Y'	X + Y	(X + Y)'	X'Y'	XY	(XY)'	X' + Y'
0	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	0	0	1	1
1	0	0	1	1	0	0	0	1	1
1	1	0	0	1	0	0	1	0	0

Generalized for n-variables

The complement of the products is the sum of the complements

The complement of the sum is the product of the complements

Ex.  $[(A' + B) C']'$

Ex.  $[(AB' + C)D' + E]'$

*Note: The complement operation is only applied to single variables in the final expression*

Dual: *Boolean expression found by replacing, AND with OR, OR with AND, 0 with 1, and 1 with 0. Variables and complements are left unchanged*

Ex.  $(XYZ)^D = X + Y + Z$        $(X + Y + Z)^D = XYZ$

Dual can be found by complementing the entire expression then complementing each individual variable

Ex.  $(AB' + C)^D$

# Laws and Theorems of Boolean Algebra

Operations with 0 and 1:

1.  $X + 0 = X$

2.  $X + 1 = 1$

1D.  $X \cdot 1 = X$

2D.  $X \cdot 0 = 0$

Idempotent laws:

3.  $X + X = X$

3D.  $X \cdot X = X$

Involution law:

4.  $(X')' = X$

Laws of complementarity:

5.  $X + X' = 1$

5D.  $X \cdot X' = 0$

Commutative laws:

6.  $X + Y = Y + X$

6D.  $XY = YX$

Associative laws:

7.  $(X + Y) + Z = X + (Y + Z)$   
 $= X + Y + Z$

7D.  $(XY)Z = X(YZ) = XYZ$

Distributive laws:

8.  $X(Y + Z) = XY + XZ$

8D.  $X + YZ = (X + Y)(X + Z)$

Simplification theorems:

9.  $XY + XY' = X$

10.  $X + XY = X$

11.  $(X + Y')Y = XY$

9D.  $(X + Y)(X + Y') = X$

10D.  $X(X + Y) = X$

11D.  $XY' + Y = X + Y$

DeMorgan's laws:

12.  $(X + Y + Z + \dots)' = X'Y'Z' \dots$

12D.  $(XYZ \dots)' = X' + Y' + Z' + \dots$

Duality:

13.  $(X + Y + Z + \dots)^D = XYZ \dots$

13D.  $(XYZ \dots)^D = X + Y + Z + \dots$

Theorem for multiplying out and factoring:

14.  $(X + Y)(X' + Z) = XZ + X'Y$

14D.  $XY + X'Z = (X + Z)(X' + Y)$

Consensus theorem:

15.  $XY + YZ + X'Z = XY + X'Z$

15D.  $(X + Y)(Y + Z)(X' + Z)$   
 $= (X + Y)(X' + Z)$

Roth, JR., C. H., & Kinney, L. L. (2010). *Fundamentals of logic design*. (Sixth ed., p. 55). Stamford, CT: Cengage.